

Control Dependence

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Chapter 7, Optimizing Compilers for Modern Architectures, Allen and Kennedy
www.cs.rice.edu/~ken/comp515/lectures/

Dependence due to Control Flow

- An example loop
 - What does the loop body (S3) depend on?

S1: if (B)
 S2: while (C)
 S3: some work

S1: if (B)
 S2: repeat
 S3: some work
 S4: until C

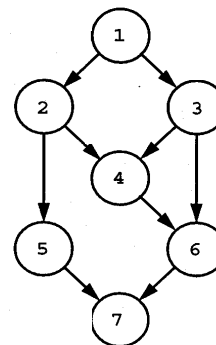
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Control Dependence

- Two definitions: A statement y control dependent on statement x if:
 - Dominance-based definition (for a compiler)
 - there exists a non-trivial path from x to y such that every statement $z \neq x$ in the path is postdominated by y and
 - x is not postdominated by y .
 - Intuitive definition (for human)
 - one branch out of x forces execution of y and another doesn't
- A node x postdominates node y if any path from y to the exit must pass through x .
- Usually we compute control dependences between basic blocks, not statements.

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What Each Node Control Dependent On?

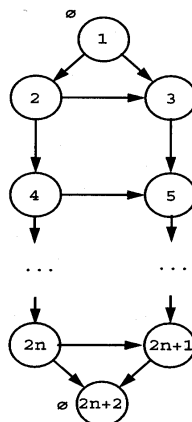


Which nodes can run in parallel, assuming no data dependence?

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Worst Case Example

- n nodes and how many control dependences?
- Can control dependence graphs get much larger than the corresponding CFG?



Control Dependence and Parallelization

- For simplicity, we shall only consider:
 - Forward branches - they create loop-independent control dependences
 - Control Dependences due to loops
- From Chapter 2: Most loop transformations are unaffected by loop-independent dependences
- Loop reversal, loop skewing, strip mining, index-set splitting, loop interchange do not affect independent dependences
- Might be problematic: Loop fusion, loop distribution
- However, since exit branches are excluded, loop fusion is not a problem

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Loop Distribution

```

DO I = 1, N
  S1   IF (A(I).LT.B(I)) GOTO 20
  S2   B(I) = B(I) + C(I)
20    CONTINUE
ENDDO

```

$S_1 \delta^{-1} S_2$

- Distributing...

```

DO I = 1, N
  S1   IF (A(I).LT.B(I)) GOTO 20
  ENDDO
DO I = 1, N
  S2   B(I) = B(I) + C(I)
  ENDDO
20    CONTINUE

```

- Incorrect!

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Loop Distribution

- Problem: control dependences crossing between distributed loops
- Solution: Keep a history of the evaluated conditions (similar to if-conversion).

```

DO I = 1, N
  S1   IF (A(I).LT.B(I)) GOTO 20
  S2   B(I) = B(I) + C(I)
20    CONTINUE
ENDDO

```

- Convert to:

```

DO I = 1, N
  S1   e(I) = A(I).LT.B(I)
  ENDDO
DO I = 1, N
  S2   IF (e(I).EQ..FALSE.) B(I) = B(I) + C(I)
  ENDDO

```

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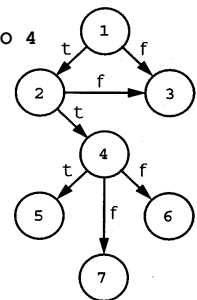
Loop Distribution

- More complex example:

```

1 DO I = 1, N
2   IF (A(I).NE.0) THEN
3     IF (B(I)/A(I).GT.1) GOTO 4
4     ENDF
5     A(I) = B(I)
6     GOTO 8
7     IF (A(I).GT.T) THEN
8       T = (B(I) - A(I)) + T
9     ELSE
10      T = (T + B(I)) - A(I)
11      B(I) = A(I)
12    ENDF
13    C(I) = B(I) + C(I)
14  ENDDO

```

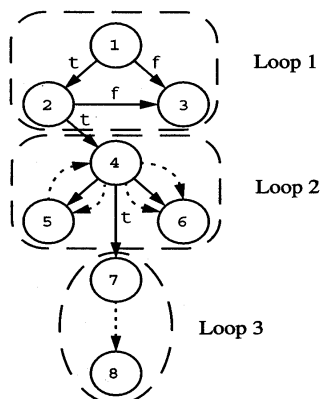


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Loop Distribution

- Fusion into "like" regions
- Needs two execution variables E2(I) and E4(I) to hold result of branches at statement 2 and 4 respectively



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